
On Robustness and Regularization of Structural Support Vector Machines: Supplementary Material

MohamadAli Torkamani
Daniel Lowd

ALI@CS.UOREGON.EDU
LOWD@CS.UOREGON.EDU

Computer and Information Science Department, University of Oregon

Proof of Lemma 4.1:

Proof. We form $\delta_{\mathcal{C}}^{\mathcal{C}}(\mathbf{x}, \mathbf{y}, \tilde{\mathbf{x}})$ (from Eq. (7) in the paper):

$$\begin{aligned} \delta^{\mathcal{C}} &= \delta_{\mathcal{C}}^{\mathcal{C}}(\mathbf{x}, \mathbf{y}, \tilde{\mathbf{x}}) = \\ \phi_{\mathcal{C}}(\tilde{\mathbf{x}}, \tilde{\mathbf{y}}) - \phi_{\mathcal{C}}(\tilde{\mathbf{x}}, \mathbf{y}) - \phi_{\mathcal{C}}(\mathbf{x}, \tilde{\mathbf{y}}) - \phi_{\mathcal{C}}(\mathbf{x}, \mathbf{y}) &= \\ \sum_{(c_x, c_y) \in \mathcal{C}} \left(\prod_{i \in c_x} \tilde{\mathbf{x}}_i - \prod_{i \in c_x} \mathbf{x}_i \right) \left(\prod_{i \in c_y} \tilde{\mathbf{y}}_i - \prod_{i \in c_y} \mathbf{y}_i \right) \end{aligned} \quad (1)$$

For an individual elements of the vector δ as expanded in (1), we can apply Hölder's inequality to the right-hand side:

$$\begin{aligned} |\delta^{\mathcal{C}}| &\leq \\ \left(\sum_{c_x \in \mathcal{C}} \left| \prod_{i \in c_x} \tilde{\mathbf{x}}_i - \prod_{i \in c_x} \mathbf{x}_i \right|^p \right)^{\frac{1}{p}} \left(\sum_{c_y \in \mathcal{C}} \left| \prod_{i \in c_y} \tilde{\mathbf{y}}_i - \prod_{i \in c_y} \mathbf{y}_i \right|^q \right)^{\frac{1}{q}} \end{aligned}$$

where $\frac{1}{p} + \frac{1}{q} = 1$. Since $\left| \prod_{i \in c_y} \tilde{\mathbf{y}}_i - \prod_{i \in c_y} \mathbf{y}_i \right|^q \leq 1$, we will have: $\sum_{c_y \in \mathcal{C}} \left| \prod_{i \in c_y} \tilde{\mathbf{y}}_i - \prod_{i \in c_y} \mathbf{y}_i \right|^q \leq |\mathcal{C}|$, therefore:

$$|\delta^{\mathcal{C}}| \leq |\mathcal{C}|^{\frac{1}{q}} \left(\sum_{c_x \in \mathcal{C}} \left| \prod_{i \in c_x} \tilde{\mathbf{x}}_i - \prod_{i \in c_x} \mathbf{x}_i \right|^p \right)^{\frac{1}{p}}$$

After applying Lemma A and raising both sides of the inequality to the power of p , we will have:

$$\begin{aligned} |\delta^{\mathcal{C}}|^p &\leq |\mathcal{C}|^{\frac{p}{q}} \left(\alpha \sum_{c_x \in \mathcal{C}} \sum_{i \in c_x} |\tilde{\mathbf{x}}_i - \mathbf{x}_i|^p \right) \\ \Rightarrow \frac{|\delta^{\mathcal{C}}|^p}{\alpha |\mathcal{C}|^{\frac{p}{q}}} &\leq \sum_{c_x \in \mathcal{C}} \sum_{i \in c_x} |\tilde{\mathbf{x}}_i - \mathbf{x}_i|^p \end{aligned} \quad (2)$$

where $\alpha = \max_{c_x \in \mathcal{C}} |c_x|^{(p-1)}$, and $|c_x|$ is the number of variables in c_x . □

The proof of Lemma 4.1 depends on the following lemma:
Lemma A. For any sequence $a_1, \dots, a_n, b_1, \dots, b_n$, such

that $0 \leq a_i, b_j \leq 1$, we have $\left| \prod_{i=1}^n a_i - \prod_{i=1}^n b_i \right|^p \leq n^{(p-1)} \sum_{i=1}^n |a_i - b_i|^p$.

Proof. For $n = 1$, the inequality is trivial. Let $u_1 = \prod_{i=1}^{\lfloor n/2 \rfloor} a_i$, $u_2 = \prod_{i=1}^{\lfloor n/2 \rfloor} b_i$, $v_1 = \prod_{i=\lfloor n/2 \rfloor + 1}^n a_i$, and $v_2 = \prod_{i=\lfloor n/2 \rfloor + 1}^n b_i$. Also it is a known fact that $|f+g|^p \leq 2^{p-1}(|f|^p + |g|^p)$ $g, f \in \mathbf{R}$. We have:

$$\begin{aligned} \left| \prod_{i=1}^n a_i - \prod_{i=1}^n b_i \right|^p &= |u_1 v_1 - u_2 v_2|^p \\ &= |u_1 v_1 - u_1 v_2 + u_1 v_2 - u_2 v_2|^p \\ &\leq 2^{p-1} (|u_1 v_1 - u_1 v_2|^p + |u_1 v_2 - u_2 v_2|^p) \\ &= 2^{p-1} (u_1^p |v_1 - v_2|^p + v_2^p |u_1 - u_2|^p) \\ &\leq 2^{p-1} (|v_1 - v_2|^p + |u_1 - u_2|^p) \end{aligned}$$

by recursive application of the above procedure, the products can be decomposed at most $\log_2 n$ times. Therefore,

$$\begin{aligned} \left| \prod_{i=1}^n a_i - \prod_{i=1}^n b_i \right|^p &\leq 2^{(p-1) \log_2 n} \sum_{i=1}^n |a_i - b_i|^p \\ &= n^{p-1} \sum_{i=1}^n |a_i - b_i|^p \end{aligned}$$

□

Proof of Corollary 4.3:

Proof. We begin with the result of Theorem 4.3, where $\frac{1}{B(d\alpha_i)^{\frac{1}{p}} |C_i|^{\frac{1}{q}}}$ is the coefficient of variations in the feature corresponding to clique C_i . Since $p = 1$ then $q = \infty$, and $\alpha_i = \max_{c_x \in C_i} |c_x|^{(p-1)} = 1$:

$$\begin{aligned} \frac{1}{B(d\alpha_i)^{\frac{1}{p}} |C_i|^{\frac{1}{q}}} &= \frac{1}{Bd |C_i|^{\frac{1}{\infty}}} \\ &= \frac{1}{Bd} \end{aligned}$$

Also in (2), set $p = 1$ and $q = \infty$.

$$|\delta^{\mathcal{C}}| \leq \left(\sum_{c_x \in \mathcal{C}} \left| \prod_{i \in c_x} \tilde{x}_i - \prod_{i \in c_x} x_i \right| \right) \max_{c_y \in \mathcal{C}} \left| \prod_{i \in c_y} \tilde{y}_i - \prod_{i \in c_y} y_i \right|$$

Since, $\max_{c_y \in \mathcal{C}} \left| \prod_{i \in c_y} \tilde{y}_i - \prod_{i \in c_y} y_i \right| = 1$, we will be using a tighter upper-bound. \square

Proof of Proposition 5.7:

Proof. We prove the case when regularization function is $\|\mathbf{w}\| = \|\mathbf{w}\|_{\infty}$ (the proofs for $\|\mathbf{M}^{-1}\mathbf{w}\|_{\infty}$ and $\|\mathbf{M}^{-1}\mathbf{w}\|_1$ are very similar, but for simplicity we chose this case). Recall that the optimization program of the robust structural SVM is:

$$\begin{aligned} & \underset{\mathbf{w}, \xi}{\text{minimize}} \quad c_1 f(\mathbf{w}) + c_2 \|\mathbf{w}\|_{\infty} + \xi \quad \text{subject to} \quad (3) \\ & \xi \geq \max_{\tilde{\mathbf{y}}} \mathbf{w}^T (\phi(\mathbf{x}, \tilde{\mathbf{y}}) - \phi(\mathbf{x}, \mathbf{y})) + \Delta(\mathbf{y}, \tilde{\mathbf{y}}) \end{aligned}$$

It can be re-written as:

$$\begin{aligned} & \underset{\mathbf{w}, \xi, t}{\text{minimize}} \quad c_1 f(\mathbf{w}) + c_2 t + \xi \quad \text{subject to} \\ & \xi \geq \max_{\tilde{\mathbf{y}}} \mathbf{w}^T (\phi(\mathbf{x}, \tilde{\mathbf{y}}) - \phi(\mathbf{x}, \mathbf{y})) + \Delta(\mathbf{y}, \tilde{\mathbf{y}}) \\ & w_i \leq t, \quad -w_i \leq t \quad \forall w_i \end{aligned}$$

In vector form we can write these constraints as: $\mathbf{w} \leq \mathbf{1}t$ and $-\mathbf{w} \leq \mathbf{1}t$. Clearly, there are two vectors \mathbf{s}_1 and \mathbf{s}_2 for which:

$$\begin{aligned} \mathbf{w} + \mathbf{s}_1 &= \mathbf{1}t \quad \Rightarrow \quad \mathbf{w} = \mathbf{1}t - \mathbf{s}_1 \\ -\mathbf{w} + \mathbf{s}_2 &= \mathbf{1}t \quad \Rightarrow \quad \mathbf{w} = \mathbf{s}_2 - \mathbf{1}t \end{aligned}$$

Let $\boldsymbol{\gamma} = [\mathbf{s}_1^T \quad \mathbf{s}_2^T \quad t]^T$, $m = \dim \mathbf{w}$, $\mathbf{I}_{\mathbf{s}_1} = [\mathbf{I}_{m \times m} \quad \mathbf{0}_{m \times m} \quad \mathbf{0}_{m \times 1}]$, $\mathbf{I}_{\mathbf{s}_2} = [\mathbf{0}_{m \times m} \quad \mathbf{I}_{m \times m} \quad \mathbf{0}_{m \times 1}]$, and $\mathbf{I}_t = [\mathbf{0}_{1 \times m} \quad \mathbf{0}_{1 \times m} \quad \mathbf{1}]$. (i.e. $\mathbf{s}_1 = \mathbf{I}_{\mathbf{s}_1} \boldsymbol{\gamma}$, $\mathbf{s}_2 = \mathbf{I}_{\mathbf{s}_2} \boldsymbol{\gamma}$, $t = \mathbf{I}_t \boldsymbol{\gamma}$). By substitution:

$$\begin{aligned} \mathbf{w} &= \mathbf{1}t \boldsymbol{\gamma} - \mathbf{I}_{\mathbf{s}_1} \boldsymbol{\gamma} = (\mathbf{1}t - \mathbf{I}_{\mathbf{s}_1}) \boldsymbol{\gamma} \\ \mathbf{w} &= \mathbf{I}_{\mathbf{s}_2} \boldsymbol{\gamma} - \mathbf{1}t \boldsymbol{\gamma} = (\mathbf{I}_{\mathbf{s}_2} - \mathbf{1}t) \boldsymbol{\gamma} \end{aligned}$$

which implies $(\mathbf{1}t - \mathbf{I}_{\mathbf{s}_1}) \boldsymbol{\gamma} = (\mathbf{I}_{\mathbf{s}_2} - \mathbf{1}t) \boldsymbol{\gamma}$, therefore: $(2 * \mathbf{1}t - \mathbf{I}_{\mathbf{s}_1} - \mathbf{I}_{\mathbf{s}_2}) \boldsymbol{\gamma} = \mathbf{0}$, or equivalently $\boldsymbol{\gamma} \in \mathcal{N}(2 * \mathbf{1}t - \mathbf{I}_{\mathbf{s}_1} - \mathbf{I}_{\mathbf{s}_2})$, where $\mathcal{N}(\cdot)$ returns the null-space of the input matrix. Let columns of matrix \mathbf{B} span $\mathcal{N}(2 * \mathbf{1}t - \mathbf{I}_{\mathbf{s}_1} - \mathbf{I}_{\mathbf{s}_2})$, also let $\boldsymbol{\gamma} = \mathbf{B}\boldsymbol{\lambda}$, we will have $\mathbf{w} = (\mathbf{1}t - \mathbf{I}_{\mathbf{s}_1}) \mathbf{B}\boldsymbol{\lambda}$. Let $\mathbf{A} = \mathbf{B}^T (\mathbf{1}t - \mathbf{I}_{\mathbf{s}_1})^T$ and $\mathbf{b} = \mathbf{I}_t^T$, then we can rewrite Problem (3) as:

$$\begin{aligned} & \underset{\boldsymbol{\lambda} \geq 0, \xi}{\text{minimize}} \quad c_1 f(\mathbf{A}^T \boldsymbol{\lambda}) + c_2 \mathbf{b}^T \boldsymbol{\lambda} + \xi \quad \text{subject to} \\ & \xi \geq \max_{\tilde{\mathbf{y}}} \boldsymbol{\lambda}^T \mathbf{A} (\phi(\mathbf{x}, \tilde{\mathbf{y}}) - \phi(\mathbf{x}, \mathbf{y})) + \Delta(\mathbf{y}, \tilde{\mathbf{y}}) \end{aligned}$$

Note that since $(2 * \mathbf{1}t - \mathbf{I}_{\mathbf{s}_1} - \mathbf{I}_{\mathbf{s}_2}) \mathbf{B} = \mathbf{0}$, we will have $(\mathbf{1}t - \mathbf{I}_{\mathbf{s}_1}) \mathbf{B} = (\mathbf{I}_{\mathbf{s}_2} - \mathbf{1}t) \mathbf{B}$, and \mathbf{A} can be transpose of any of them. \square