#### **Finite Models**



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# Properties of Models

- Compact: representable and manipulable in a reasonably compact form
  - What is *reasonably compact* depends largely on how the model will be used
- Predictive: must represent some salient characteristics of the modeled artifact well enough to distinguish between good and bad outcomes of analysis
  - no single model represents all characteristics well enough to be useful for all kinds of analysis
- Semantically meaningful: it is usually necessary to interpret analysis results in a way that permits diagnosis of the causes of failure
- Sufficiently general: models intended for analysis of some important characteristic must be general enough for practical use in the intended domain of application

# Learning objectives

- Understand goals and implications of finite state abstraction
- Learn how to model program control flow with graphs
- Learn how to model the software system structure with call graphs
- Learn how to model finite state behavior with finite state machines



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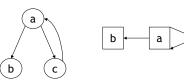
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# Graph Representations: directed graphs

- Directed graph:
  - N (set of nodes)
  - E (relation on the set of nodes ) edges

Nodes: {a, b, c} Edges: {(a,b), (a, c), (c, a)}

[(a,b), (a, c), (c, a)}

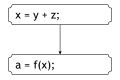




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## Graph Representations: labels and code

- We can label nodes with the names or descriptions of the entities they represent.
  - If nodes a and b represent program regions containing assignment statements, we might draw the two nodes and an edge (a,b) connecting them in this way:





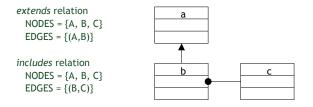
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#### Multidimensional Graph Representations

- Sometimes we draw a single diagram to represent more than one directed graph, drawing the shared nodes only once
  - class B extends (is a subclass of) class A
  - class B has a field that is an object of type C



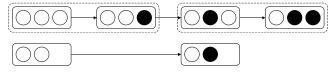


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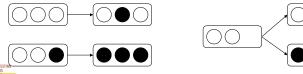
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#### Finite Abstraction of Behavior

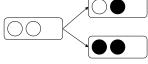
an abstraction function suppresses some details of program execution



it lumps together execution states that differ with respect to the suppressed details but are otherwise identical







### (Intraprocedural) Control Flow Graph

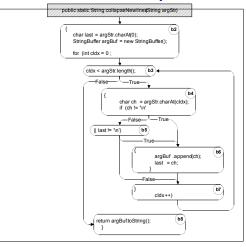
- nodes = regions of source code (basic blocks)
  - Basic block = maximal program region with a single entry and single exit point
  - Often statements are grouped in single regions to get a compact model
  - Sometime single statements are broken into more than one node to model control flow within the statement
- directed edges = possibility that program execution proceeds from the end of one region directly to the beginning of another



# **Example of Control Flow Graph**

```
public static String collapseNewlines(String argStr)
{
    char last = argStr.charAt(0);
    StringBuffer argBuf = new StringBuffer();

    for (int cldx = 0 ; cldx < argStr.length(); cldx++)
    {
        char ch = argStr.charAt(cldx);
        if (ch != "\n" || last != "\n")
        {
            argBuf.append(ch);
            last = ch;
        }
    }
    return argBuf.toString();
}</pre>
```



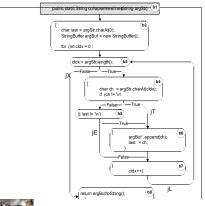


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#### Linear Code Sequence and Jump (LCSJ)

Essentially subpaths of the control flow graph from one branch to another



From	Sequence of basic blocs	То
Entry	b1 b2 b3	jΧ
Entry	b1 b2 b3 b4	jΤ
Entry	b1 b2 b3 b4 b5	jΕ
Entry	b1 b2 b3 b4 b5 b6 b7	jL
jΧ	b8	ret
jL	b3 b4	jΤ
jL	b3 b4 b5	jΕ
jL	b3 b4 b5 b6 b7	jL

3

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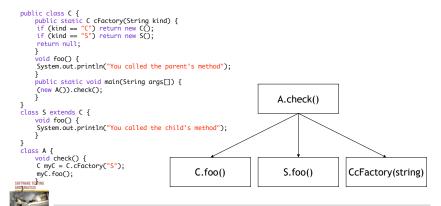
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# Interprocedural control flow graph

- Call graphs
  - Nodes represent procedures
    - Methods
    - C functions
    - ..
  - Edges represent calls relation

# Overestimating the *calls* relation

The static call graph includes calls through dynamic bindings that never occur in execution.





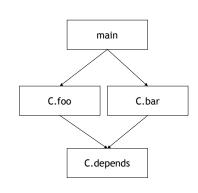
### Contex Insensitive Call graphs

```
public class Context {
  public static void main(String args[]) {
     Context c = new Context();
     c.foo(3);
     c.bar(17);
  }

  void foo(int n) {
    int[] myArray = new int[ n ];
    depends( myArray, 2);
  }

  void bar(int n) {
    int[] myArray = new int[ n ];
    depends( myArray, 16);
  }

  void depends( int[] a, int n ) {
    a[n] = 42;
  }
}
```

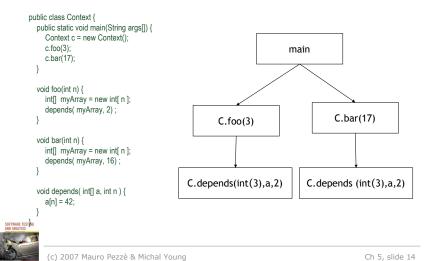




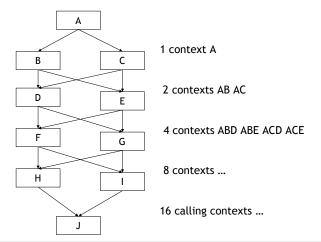
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#### Contex Sensitive Call graphs



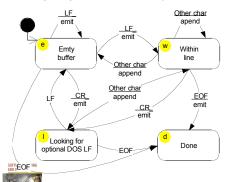
# Context Sensitive CFG exponential growth



#### Finite state machines

- · finite set of states (nodes)
- · set of transitions among states (edges)

#### Graph representation (Mealy machine)



#### Tabular representation

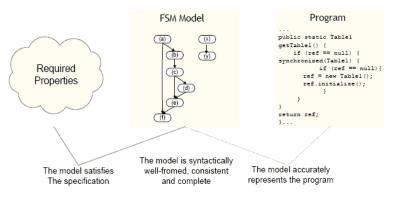
	LF	CR	EOF	other
е	e/emit	e/emit	d/-	w/append
w	e/emit	e/emit	d/emit	w/append
I	e/-		d/-	w/append

AND ANALYSIS

(C)

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# Using Models to Reason about System Properties





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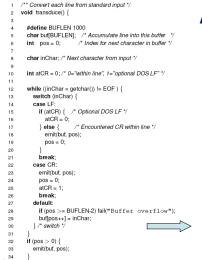
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# Summary

- Models must be much simpler than the artifact they describe to be understandable and analyzable
- Must also be sufficiently detailed to be useful
- CFG are built from software
- FSM can be built before software to documentintended behavior



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#### **Abstraction Function**

Abstract state	Concrete state		
	Lines	atCR	pos
e (Empty buffer) w (Within line) l (Looking for LF)	3 - 13	0	0
w (Within line)	13	0	> 0
l (Looking for LF)	13	1	0
d (Done)	36	-	-



	LF	CR	EOF	other
e	e / emit	1/emit	d/-	w / append
W	e / emit	1/ emit	d/emit	w / append
1	e/-	1/emit	d/-	w / append



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